

Research Proposal: A data structure to support the simulation of random events

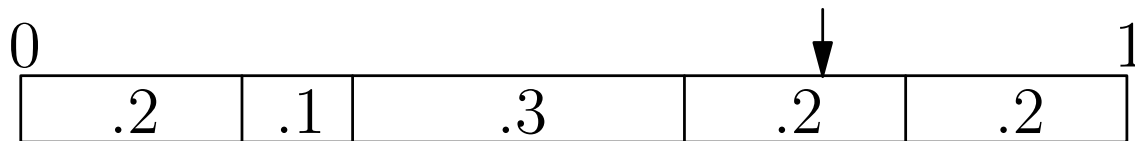
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Outline

- the random object selection problem
- reaction-diffusion equation (motivation)
- problem description
- related problems:
 - selectable partial sums
 - optimal search trees/coding systems
- research goals and approaches

Random object selection problem

- Given:
 - n objects
 - relative probabilities p_1, \dots, p_n ($\sum_{i=1}^n p_i = 1$)
- Goal: select an object at random
 - the probability of selecting object k is p_k
- Static solution:
 - precompute $\sigma_k = \sum_{i=1}^k p_i$
 - generate uniform- $[0, 1)$ random variable x
 - binary search for k s.t. $\sigma_{k-1} \leq x < \sigma_k$



$$\sigma_k: \boxed{0 \mid .2 \mid .3 \mid .6 \mid .8 \mid 1}$$

Random object selection problem

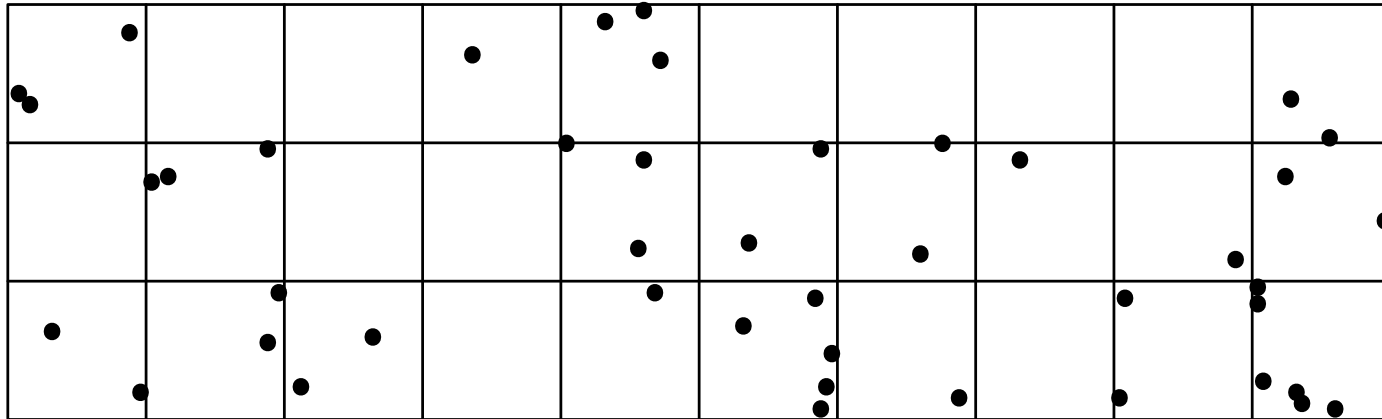
- limitations:
 - probabilities cannot change
 - set of objects cannot change
 - all objects treated equally

Reaction-diffusion equation

- describes how a chemical (or other contaminant) behaves in a fluid
 - chemical movement (e.g. diffusion, fluid flow)
 - increases/decreases in concentration (e.g. reactions, births/deaths)

Simulating

- divide fluid into cells
- concentration \rightarrow discrete number of particles
 - movement of chemical \rightarrow movement of particles between adjacent cells
 - increases/decreases in concentration \rightarrow particle “births” and “deaths”



Simulating

state: collection of all the particle counts

event: e , a particle movement, birth, or death

- modeled as Markov process
- time of next occurrence: exponential random variable with rate λ_e
- function of particle counts, spatial location

Straightforward simulation

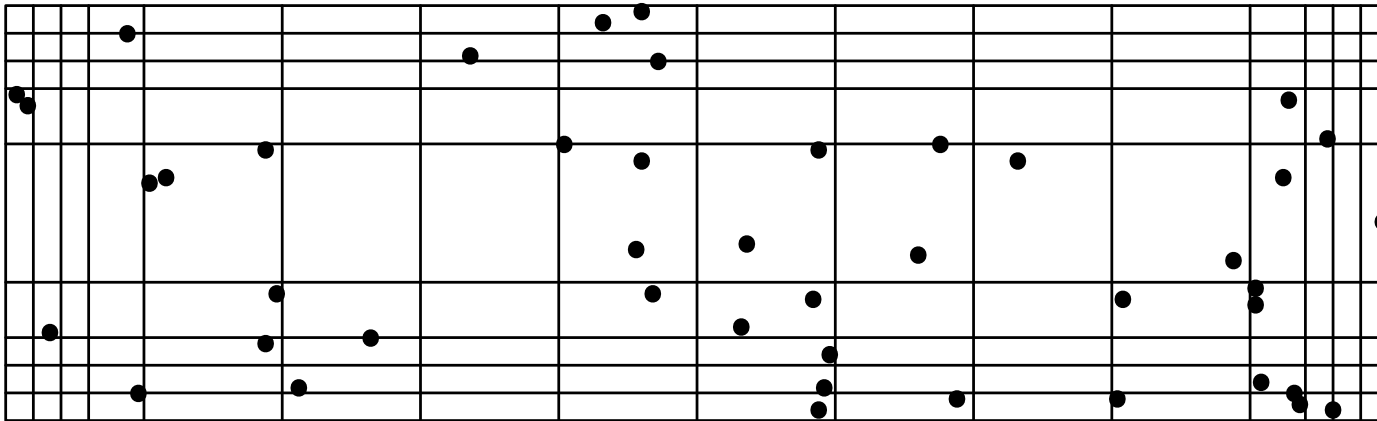
- generate n exponential random variables (time of events)
- find random variable. with smallest value (select event)
- simulate event
- update rates
- re-generate random variables as needed, due to changed rates

More efficient simulation

- let $\lambda_E = \sum_{e \in E} \lambda_e$, sum of all rates
- generate one λ_E -exponential r.v. (time of events)
- select event randomly s.t. event e has probability
 $p_e = \lambda_e / \lambda_E$
- simulate event
- update rates

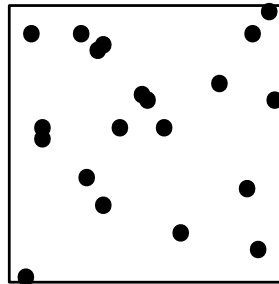
Splitting and merging

- some areas are of greater interest
 - e.g. modeling pollution in a lake → more interested in areas close to shore
 - grid may be non-uniform



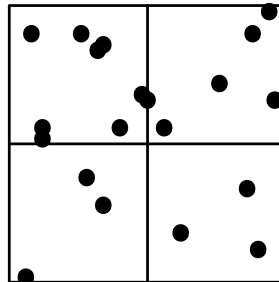
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- grid may change dynamically
 - e.g. more interested in areas of high concentration
 - split/merge cells



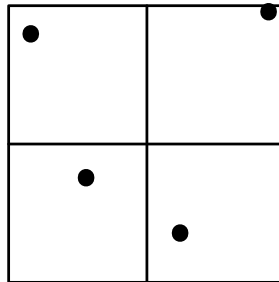
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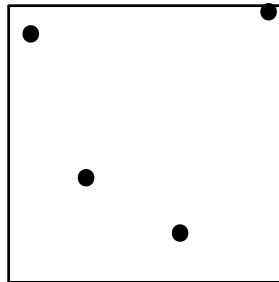
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Problem description

- create a data structure
- Given: set of pairs: events and their rates
 $(e_1, \lambda_1), \dots, (e_n, \lambda_n)$
- support these operations:
 - random select:** select an event at random, with probability λ_e / λ_E
 - update:** set the rate of an event to a new value
 - split:** split one event \rightarrow a number of events, distribute rate evenly
 - merge:** merge a number of events \rightarrow one event with rate equal to sum

Problem description

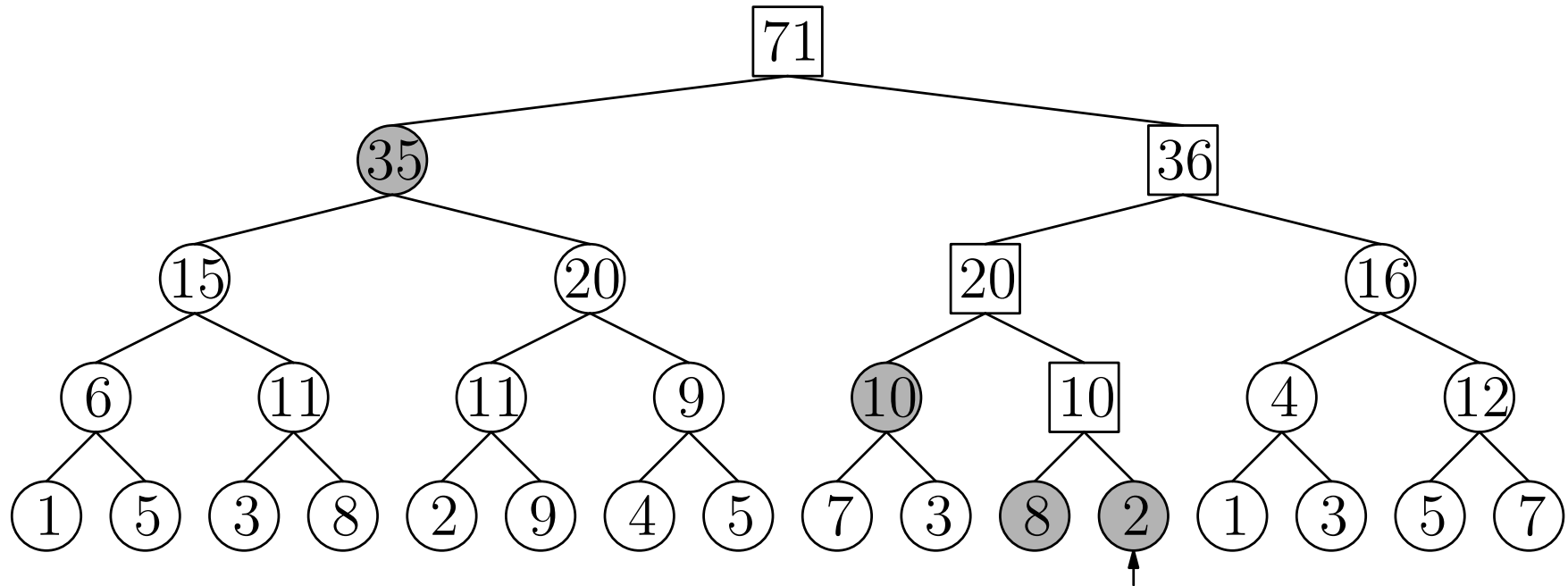
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 - merge:** merge a number of events \rightarrow one event with rate equal to sum
 - insert:** add a new event along with its rate
 - delete:** remove an event along with its rate

Selectable partial sums

- related to simple random object selection problem
- given non-negative *keys* a_1, \dots, a_n
- support these operations:
 - update:** set the value of a key to a new value
 - sum:** given k , calculate $\sigma_k = \sum_{i=1}^k a_i$
 - select:** given target t , find k such that $\sigma_{k-1} \leq t < \sigma_k$
- to solve random object selection problem:
 - probabilities $p_i \leftrightarrow a_i$
 - uniform-[0, 1) random variable $x \leftrightarrow t$

Binary trees and selectable partial sums

- simplest sub-linear data structure



Previous work on selectable partial sums

Pătrașcu and Demaine: $\Theta(1 + \lg n / \lg(b/\delta))$ on b -bit machine, δ -bit additive changes (upper and lower bounds)

Raman, Raman, Rao: *succinct data structure* $kn + o(kn)$ space, $O(\lg n / \lg \lg n)$ time

Hon, Sadakane, Sung: keys up to $O(\lg \lg n)$ bits, trade off between update and queries

Moffat: $O(\log(1 + k))$ time to update, sum, or select k th key a_k

Hampapuram and Fredman: updates and sums have different probabilities, but does not support selections

Optimal search trees/coding systems

- use a tree to solve selectable partial sums
- consider different access probabilities
- optimal search trees: minimize expected cost for a search
 - entropy $H = \sum_{i=1}^n p_i \log 1/p_i$
 - expected search time: between $H - \log H - \log e + 1$ and $H + 2$
- some coding systems (e.g. Huffman): correspondence with trees
 - also tries to minimize expected access cost

Previous work on coding

Faller, Gallager, Knuth: dynamic Huffman coding

Vitter: improved FGK algorithm

Gagie: dynamic Shannon coding

Research goals

- support the operations with the following running times:

select: $O(\log 1/p)$

update: $O(\log 1/p)$, or maybe $O(f(\Delta\lambda) \log 1/p)$

split: ?

merge: ?

insert: $O(\log 1/p)$

delete: $O(\log 1/p)$

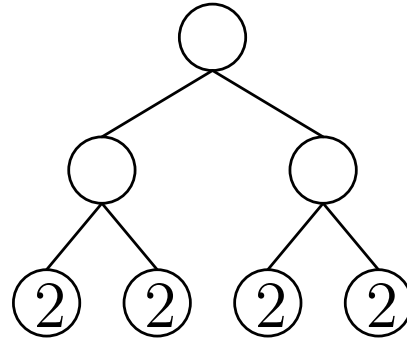
Splitting

- under a standard tree: trivial

⑧

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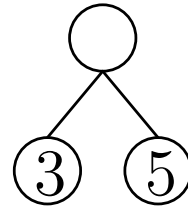


Splitting

- under a standard tree: trivial
- harder if tree is stored in special structure (e.g. Vitter), or data structure has other constraints

Merging

- harder than splitting
- if nodes to be merged have common parent → easy



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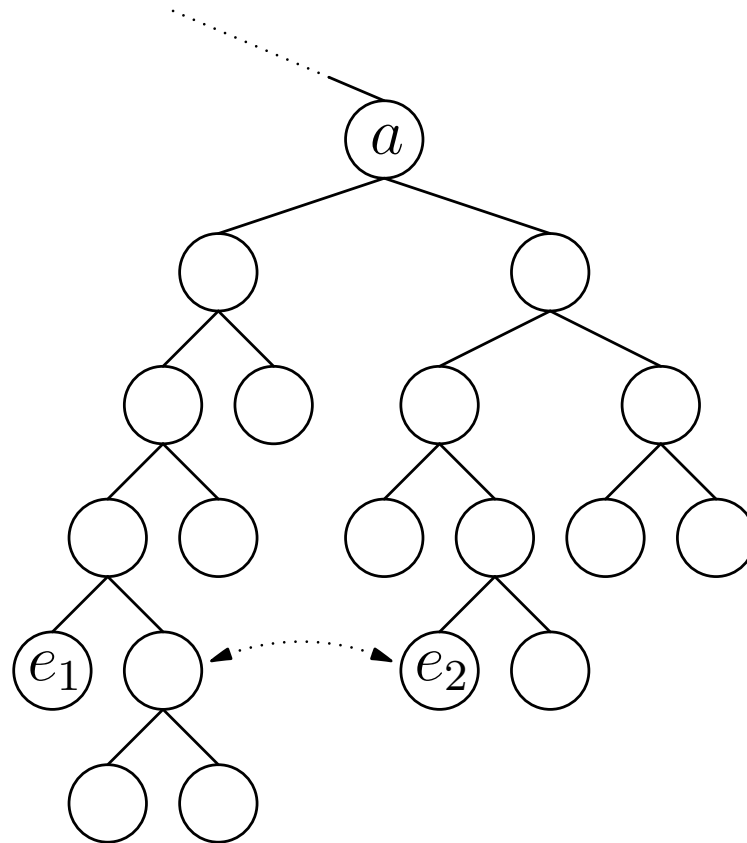
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Merging

- harder than splitting
- if nodes to be merged have common parent → easy
- otherwise, need to remove old nodes, insert new node
→ minimize time

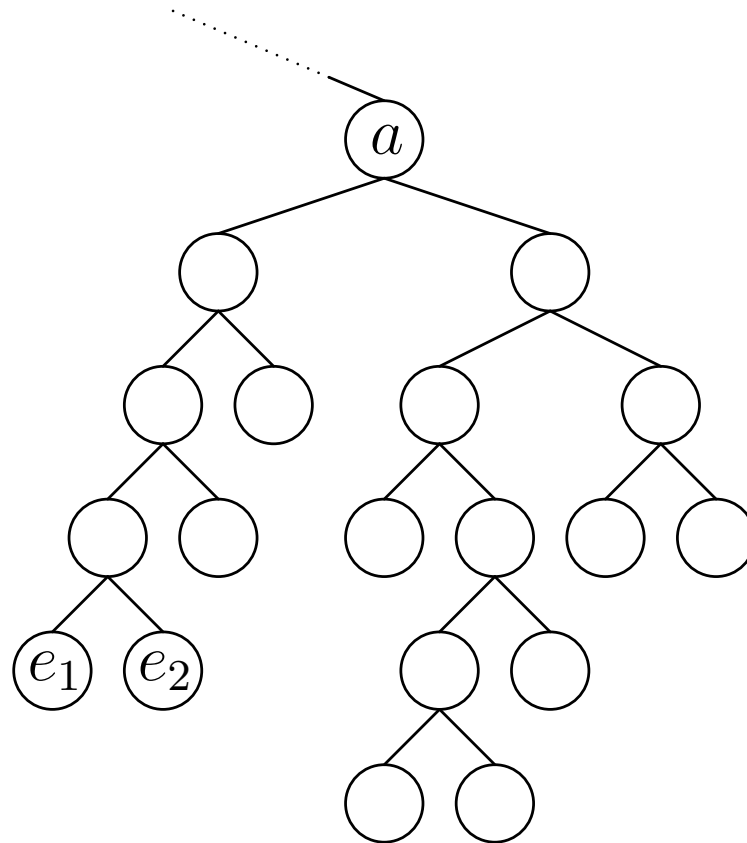
Merging

- if events on same level:



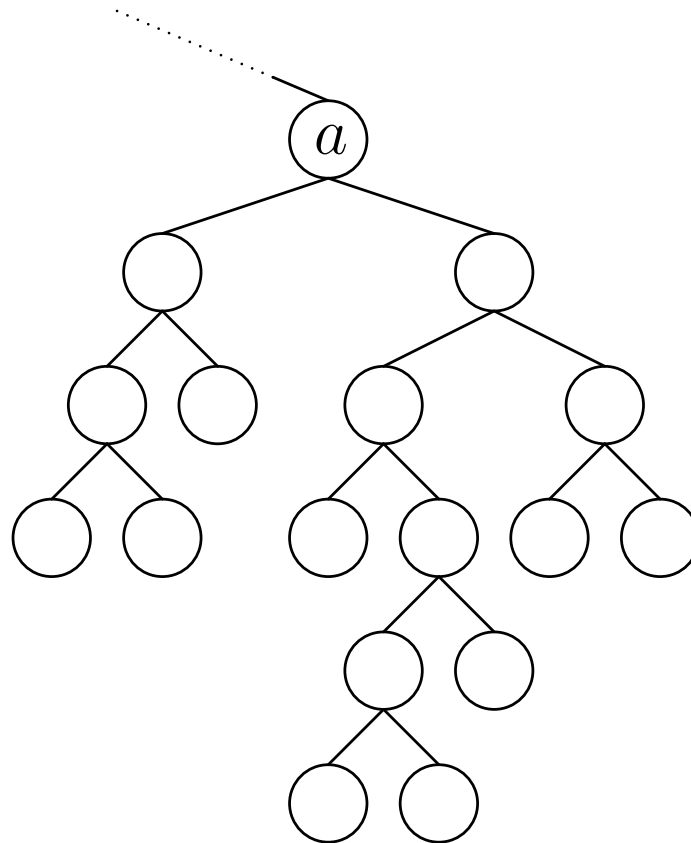
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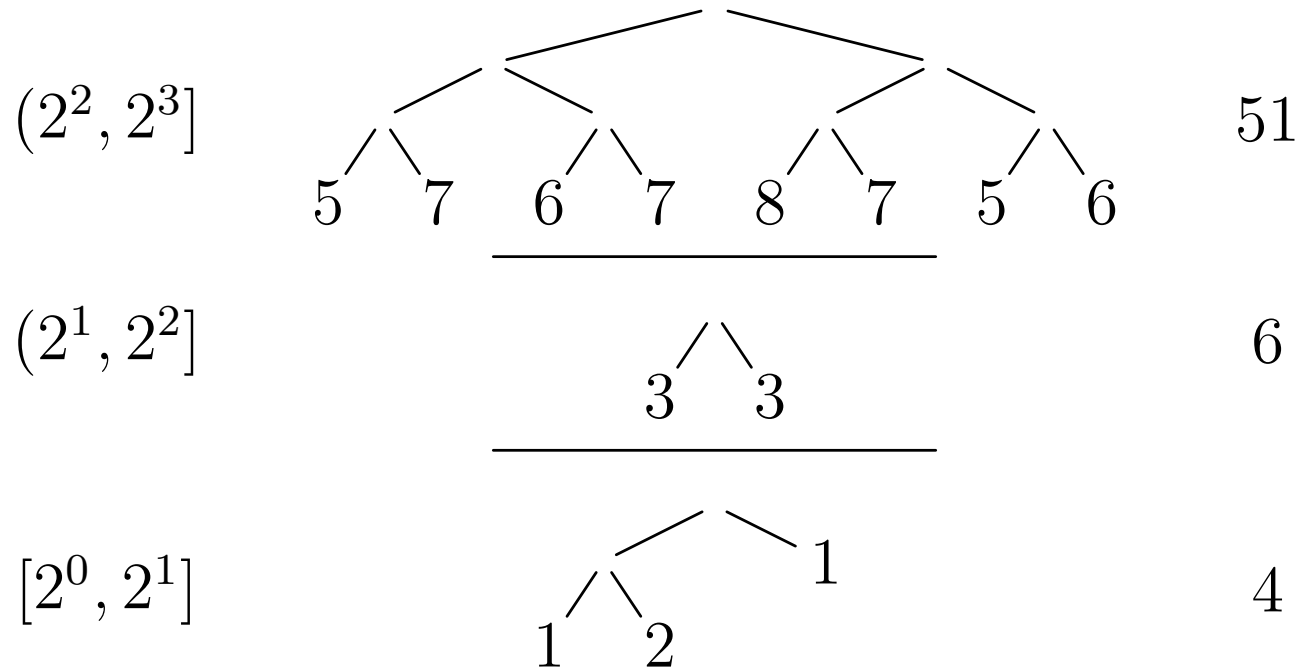


Approaches

- try to adapt previous work
 - Hampapuram and Fredman: consider different probabilities, but do not support needed operations
 - coding systems
 - support updating weights, but only increment/decrement by one $\rightarrow O(\Delta\lambda \log 1/p)$
 - also can support insert/delete, but too slow

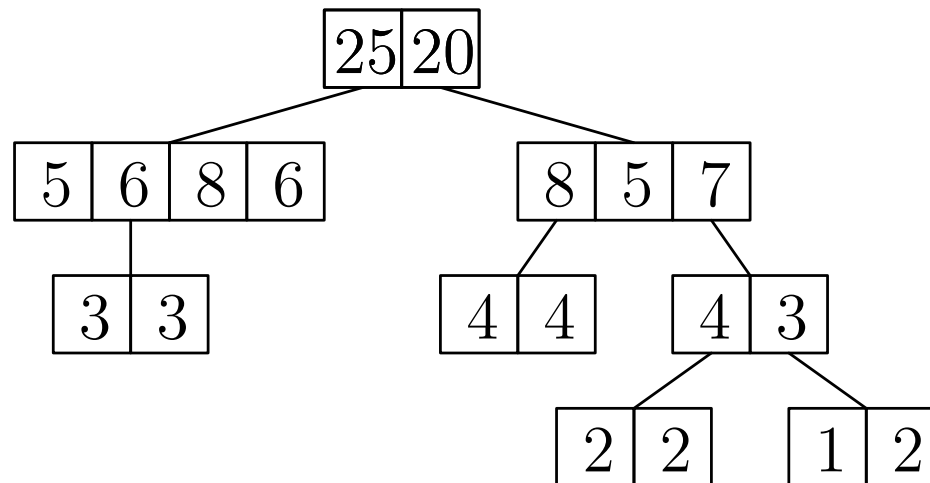
Approaches

- group nodes according to rate
- put each group in a balanced tree



Approaches

- B-tree-like structure
- elements in each node are within a range of rates



Approaches

- divide space into cells
- similar to quad trees
- need further research

Summary

- random object selection problem with split and merge
- speed up simulation of reaction-diffusion equation
- related to partial sums, optimal search trees, coding systems
 - starting point in our research
- other approaches based on
 - grouping nodes by rates
 - B-trees
- quad trees